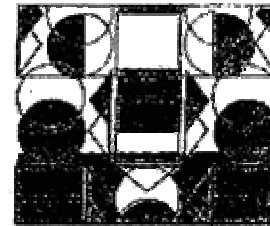
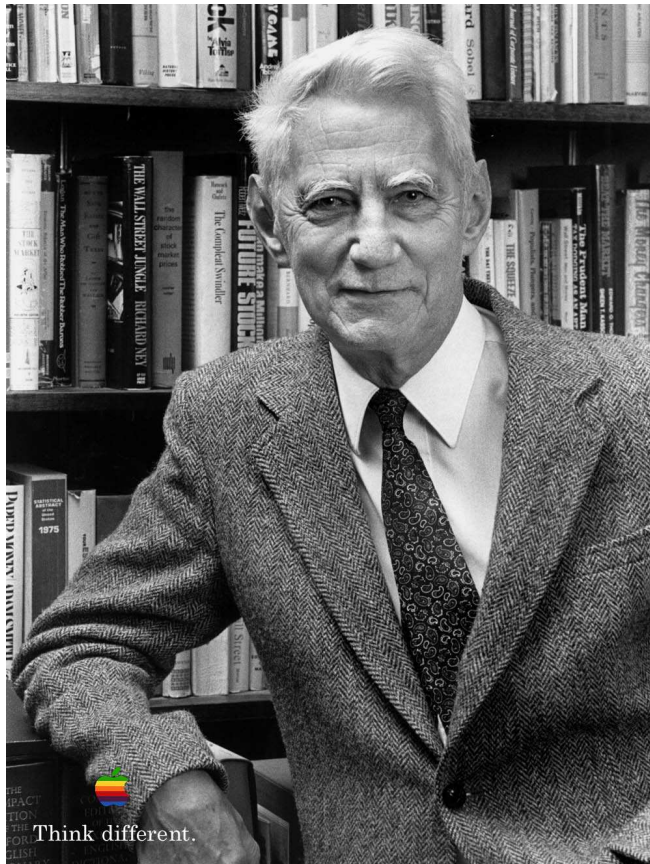

Information Theory

Laura Kovács



THE MATHEMATICAL THEORY OF COMMUNICATION

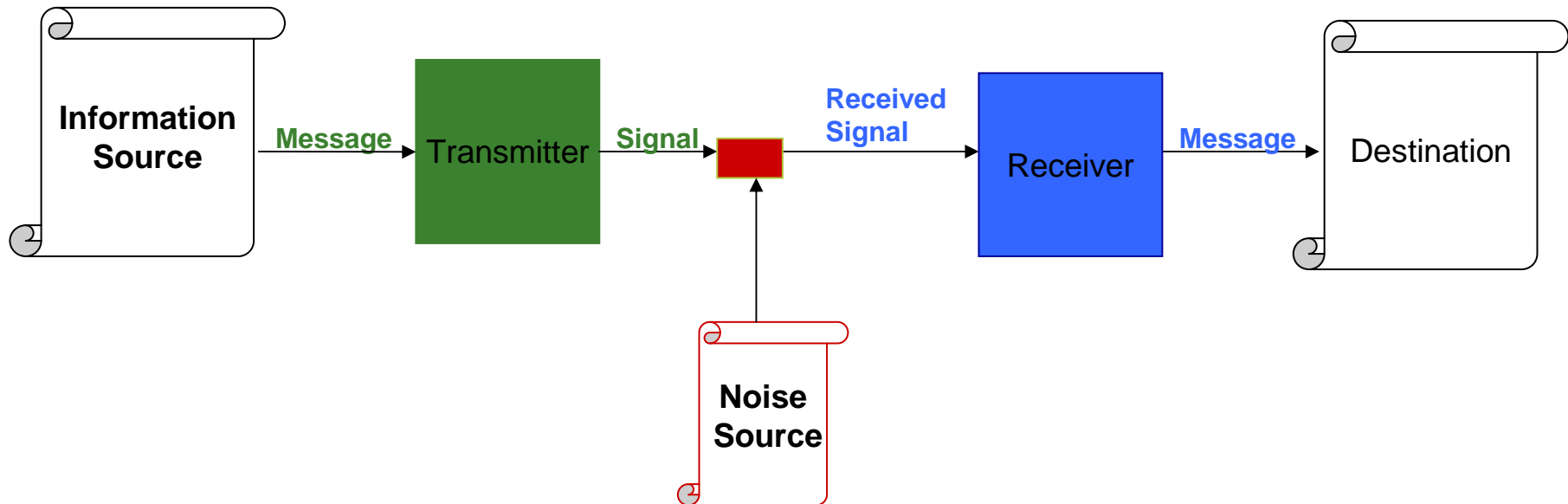
by Claude E. Shannon and Warren Weaver

1949

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Claude Shannon (1916-2001)

A Communication System Model



Information Theory

by Claude Shannon & Warren Weaver

- **Information** (measured in *bits*):
 - the amount of uncertainty a message eliminates;
 - Based on a **linear model** of communication;
 - **Noise** :
 - reduces information by increasing uncertainty;
 - **Redundancy**:
 - reduces loss of information due to noise.
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e an i 0 du an
Ge an is 0% red nt
German is 50% redunda

Information Content $h(p)$

(Informationgehalt)

= **MEASURE OF UNCERTAINTY** associated with probability p

- Independent of coding;
 - Increases when probability p decreases (uncertainty increases);
-

Information Content $h(p)$

(Informationgehalt)

= **MEASURE OF UNCERTAINTY** associated with probability p

- Independent of coding;
 - Increases when probability p decreases (uncertainty increases);
 - $h(p_1 * p_2) = h(p_1) + h(p_2)$;
 - $h(p) = \text{ld}(1/p)$ [bit].
-

Examples of Information Content

- $h(1) = 0$ Message will be always expected
 - $h(0) = \infty$ Message will be never expected
 - $h(0.5) = 1$ Message will be by 50% expected
 - $h(0.1) = 3.32$ Information content of a decimal digit
-

Binary Logarithm $\text{ld } x$

Properties, identities

- $2^{\text{ld } x} = x$
 - $\text{ld } 1/x = -\text{ld } x$
 - $\text{ld } (x*y) = \text{ld } x + \text{ld } y$
 - $\text{ld } 2*x = 1 + \text{ld } x$
 - $\text{ld } x^n = n*\text{ld } x$
 - $\text{ld } x^2 = 2*\text{ld } x$
 - $\text{ld } x = \ln x / \ln 2 \sim 1.44*\ln x$
 - $\text{ld } x = \log x / \log 2 \sim 3.32*\log x$
-

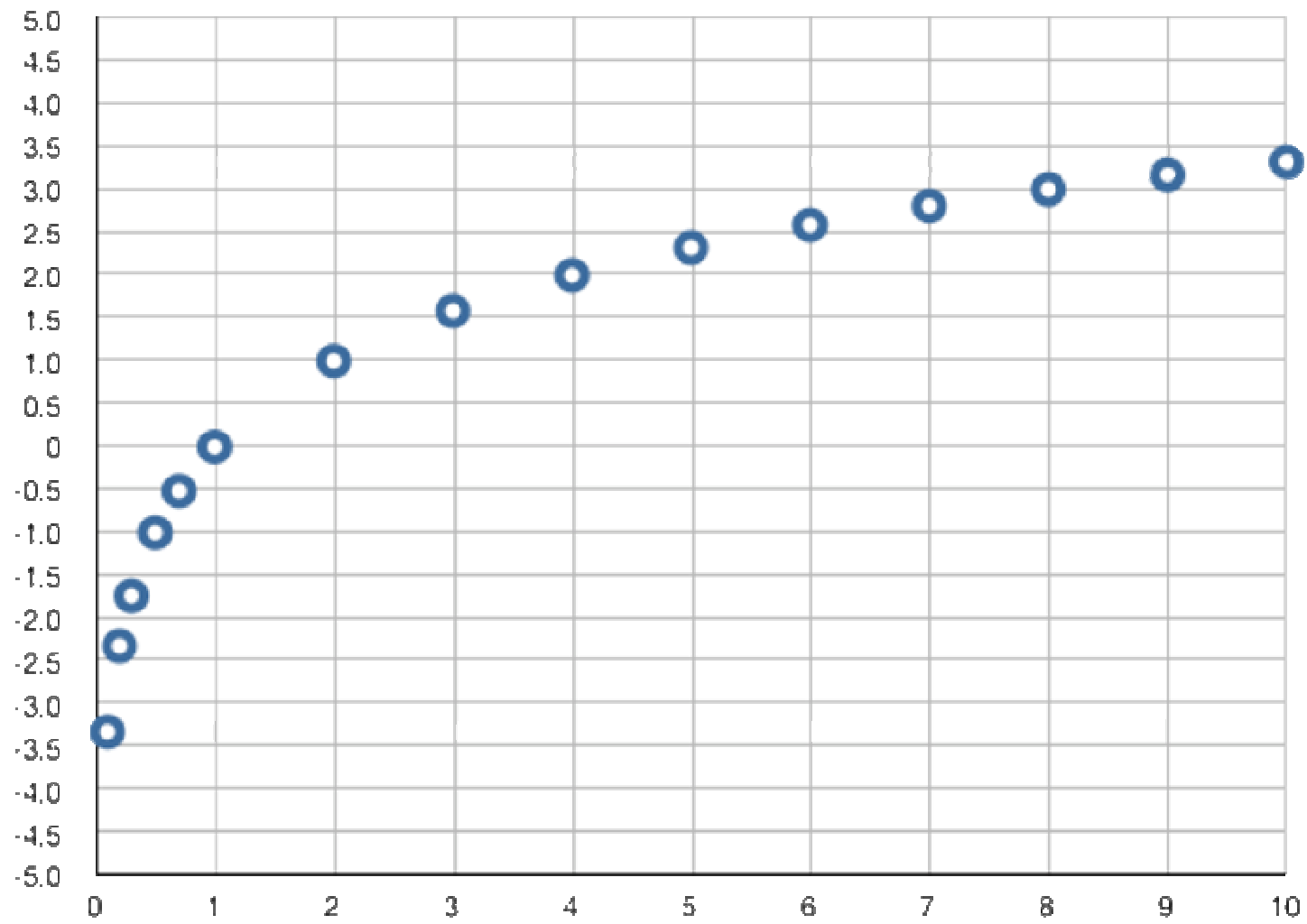
Examples

- $\text{ld } 1 = 0$
- $\text{ld } 2 = 1$
- $\text{ld } 4 = \text{ld } (2^2) = 2 * \text{ld } 2 = 2$
- $\text{ld } 8 = \text{ld } (2^3) = 3 * \text{ld } 2 = 3$
- $\text{ld } 10 \sim 3.32 * \log 10 = 3.32 * 1 = 3.32$
- $\text{ld } 5 = \text{ld } 10 - 1 \sim 2.32$
- $\text{ld } 9 \sim (\text{ld } 8 + \text{ld } 10) / 2 = (3 + 3.32) / 2 = 3.16$
- $\text{ld } 3 = (\text{ld } 9) / 2 \sim 1.58$
- $\text{ld } 7 = (\text{ld } 49) / 2 \sim (\text{ld } 50) / 2 = (\text{ld } 100 - 1) / 2 = (2 * \text{ld } 10 - 1) / 2 = 2.82$

x

ld x

1	0.00
2	1.00
3	1.58
4	2.00
5	2.32
6	2.58
7	2.81
8	3.00
9	3.17
10	3.32



- **Average Information Content (Entropy):**

[Mittlerer Informationsgehalt]

$$H = \sum_i p_i \log 1/p_i \text{ [bit]}$$

- **Average Word Length:**

[Mittlere Wortlänge]

$$L = \sum_i p_i l_i \text{ [bit]}$$

- **Redundancy:**

[Redundanz]

$$R = L - H \text{ [bit]}$$

Shannon's Coding Theorem: $L \geq H$

Shannon's Binary Entropy Function

[Shannon'sche Binäre Funktion]

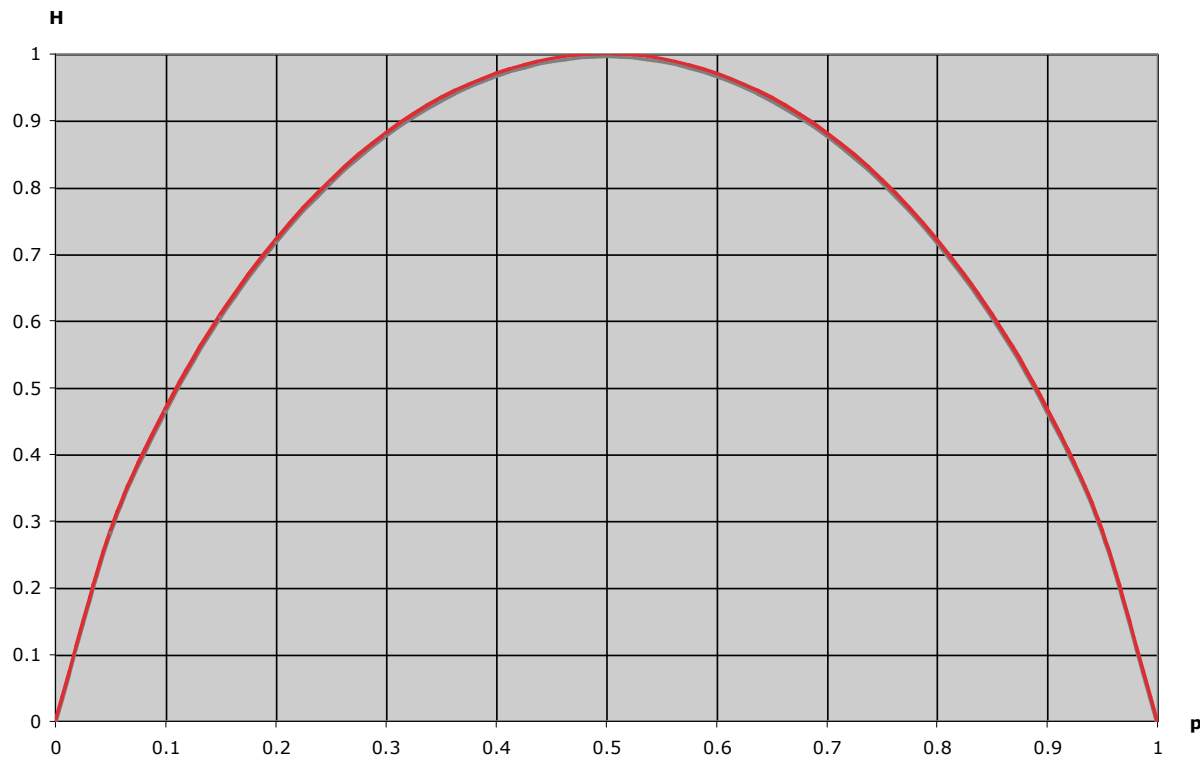
Example 1:

Consider tossing a coin with known, not necessarily fair, probabilities of coming up heads or tails.

Shannon's Binary Entropy Function

[Shannon'sche Binäre Funktion]

$$H = -p \log_2 1/p - (1-p) \log_2 1/(1-p)$$



Entropy of a coin toss as a function of the probability p of it coming up heads.

Example 2: Redundancy of a decimal digit.

(All 10 digits have same probabilities.)

- Average Information Content (Entropy):

$$H = \log_2 10 = 3.32 \text{ bit}$$

- Word Length (each binary digit has the *same prefixed length* of 4 bits):

$$L = 4 \text{ bit}$$

- Redundancy:

$$R = L - H = 0.68 \text{ bit}$$

Example 3: Average Information Content of a (German) Letter.
(All 26 letters have same probabilities.)

- Average Information Content (Entropy):

$$H = \log_2 26 = 4.7 \text{ bit}$$



Example 4: Average Information Content of a (German) Letter.
(Letters have different probabilities.)

- Average Information Content (Entropy):

$$H = 4.1 \text{ bit}$$

Note: More than a half of all German texts are formed by letters e, n, r, i and the space symbol.

Example 4: Average Information Content of a (German) Letter.
(Sequence of [2] consecutive letters have different probabilities.)

- Average Information Content (Entropy):

H ~ 2 bit

Example 5: Average Information Content of a (German) Word.
(10 Million Words have different probabilities.)

- Average Information Content (Entropy):

$$H \sim 15.5 \text{ bit}$$

- Word Length:

$$L = 7.7 \text{ letters}$$

Example 6: Information flow during reading.

Reading: 25 letters per second



Information flow: 50 bits/sec



Storing: 25 bits/sec



In 60 years a person could store approx. $3 \cdot 10^{10}$ bits.



Theoretically: the brain's storage capacity ~ 10^{12} bits.

→ one could store read information for more than ~1900 years 😊

Reminder for Next Lecture

- Carry out the solution of Example 4 of this lecture!
 - Next lecture on October 7th !
 - No lecture on September 30th .
 - Homework sessions start today (September 23rd), at 12pm!
 - This week's homework due by September 30th, 10am.
 - Contact teaching assistants for help, if needed!
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